Understanding Operating Systems 5th Edition McHoes Test Bank

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Chapter 2: Memory Management: Early Systems

TRUE/FALSE

1.	Early memory management schemes are still used in today's operating systems.						
	ANS:	F	PTS:	1	REF:	32	
2.						ry and if a program doesn't fit, then either the size of m must be modified.	
	ANS:	T	PTS:	1	REF:	32	
3.						om main memory into secondary storage for s occupying the same memory locations.	
	ANS:	F	PTS:	1	REF:	32	
4.				b in a single-us emory protection		em is storing the first memory location of program	
	ANS:	T	PTS:	1	REF:	33	
5.	A sing	gle-user system	suppor	ts multiprogran	nming.		
	ANS:	F	PTS:	1	REF:	33	
6.	The first attempt to allow for multiprogramming used fixed partitions.						
	ANS:	T	PTS:	1	REF:	34	
7.	The pi	roblem of parti	tion intr	rusion is presen	t in sin	gle-user contiguous allocation schemes.	
	ANS:	F	PTS:	1	REF:	34	
8.	The algorithm used to store jobs into memory requires a few more steps than the one used for a single-user system because the size of the job must be matched with the size of the partition to make sure it fits completely.						
	ANS:	T	PTS:	1	REF:	34	
9.	The fixed partition scheme does not require that the entire program be stored contiguously and in memory from the beginning to the end of its execution.						
	ANS:	F	PTS:	1	REF:	35	
10.		•				jobs run on the system are of the same size or if the veen reconfigurations.	
	ANS:	T	PTS:	1	REF:	35	

11.	In a fixed partition scheme, large jobs may have a longer turnaround time as they wait for free partitions of sufficient size or may never run.					
	ANS: T	PTS:	1	REF:	36	
12.	The best-fit allocati memory to high-ord			ee/busy	lists organized by memory locations, low-order	
	ANS: F	PTS:	1	REF:	38	
13.	A large job can hav	e problei	ns with a first-	fit mem	nory allocation list.	
	ANS: T	PTS:	1	REF:	39	
14.	The best-fit free list to implement.	scheme	uses memory n	nore ef	ficiently than the first-fit free scheme but it is slower	
	ANS: T	PTS:	1	REF:	40	
15.	The first-fit algorith blocks.	nm assum	nes that the Men	mory M	Ianager keeps only one list containing free memory	
	ANS: F	PTS:	1	REF:	40	
16.	One of the problems with the best-fit algorithm is that the entire table must be searched before the allocation can be made because the memory blocks are physically stored in sequence according to their location in memory.					
	ANS: T	PTS:	1	REF:	42	
17.	Research continues	to focus	on finding the	optimu	m allocation scheme.	
	ANS: T	PTS:	1	REF:	43	
18.	For a fixed partition	ı system,	memory deallo	ocation	is quite complex.	
	ANS: F	PTS:	1	REF:	44	
19.	A null entry in the best returned to the free		occurs when a i	memory	y block between two other busy memory blocks is	
	ANS: T	PTS:	1	REF:	47	
20.		mpty blo	cks and compa	ct them	Memory Manager relocates programs to gather to make one block of memory large enough to t in.	
	ANS: T	PTS:	1	REF:	48	
21.	Memory is allocated	d during	garbage collect	ion.		
	ANS: F	PTS:	1	REF:	48 49	

22.	. During compaction, the operating system must distinguish between addresses and data values, and the distinctions are not obvious once the program has been loaded into memory.					
	ANS: T	PTS: 1	REF:	49		
23.	After relocation ar	nd compaction, both	th the free list	and the busy list are updated.		
	ANS: T	PTS: 1	REF:	51		
24.	The bounds registed in memory accession			lowest, depending on the specific system) location		
	ANS: T	PTS: 1	REF:	52		
25.	Compaction shoul	d always be perfor	med only whe	on there are jobs waiting to get in.		
	ANS: F	PTS: 1	REF:	53 54		
MUL	TIPLE CHOICE					
1.	 The following,, describes the first memory allocation scheme. a. Each program to be processed was loaded into secondary storage, then swapped into memory in parts b. Each program to be processed was partially loaded into memory, then granted more memory as needed c. Each program to be processed was allocated a portion of memory and could negotiate with other programs to access more memory d. Each program to be processed was loaded in its entirety into memory and allocated as much contiguous space in memory as it needed 					
	ANS: D	PTS: 1	REF:	32		
2.	In a single-user sy a. sequentially b. intermittently	stem, jobs are prod	c.	randomly in order of longest job to shortest job		
	ANS: A	PTS: 1	REF:	32		
3.	 In the algorithm to load a job in a single-user system, the program counter is initially set to a. the address of the last memory location b. the number of instructions c. zero d. the address of the first memory location 					
	ANS: D	PTS: 1	REF:	33		
4.	Fixed partitions ar a. complete b. static	re also called	c.	direct sized		
	ANS: B	PTS: 1	REF:	34		
5.	a. Comparing job			in a fixed partition.		

b. Determining the job's requested memory sizec. Setting counter to oned. Placing the job in a waiting queue						
	AN	IS: B	PTS:	1	REF:	34
6.	the a. b. c.	partition scher partition size, mostatus, access, ar partition size, sta partition size, mo	emory and memo	ddress, and star ory address I access	tus	Manager uses to keep track of jobs is composed of tus
	AN	IS: D	PTS:	1	REF:	35
7.	a. b. c.	when jobs have when jobs have when jobs sizes a when all jobs are	the same differen re not k	e size t sizes nown in advand		
	AN	IS: A	PTS:	1	REF:	35
8.	a.		t	of free memory	c.	n blocks of allocated memory. External fragmentation Internal fragmentation
	AN	IS: C	PTS:	1	REF:	36
9.	me a. b.	e keeps the famory. fixed partition al first-fit memory IS: B	location	n on	c.	dynamic fit memory allocation best-fit memory allocation
10.		has the least was Fixed partitioning. First-fit memory	ıg		•	artition fitting the requirements. Dynamic fit memory allocation Best-fit memory allocation
	AN	IS: D	PTS:	1	REF:	38
11.	sch Job J1 J2 J3 J4 Blo B1 B2 B3 B4	neme is used, the jos: 10K 20K 30K 10K ccks: 30K 15K 50K			run.	s 1-4 and memory blocks. Assuming a first-fit
	a.	Jl			C.	J3

```
b. J2
                                                  d. J4
    ANS: C
                         PTS: 1
                                              REF: 39|40
12. Consider the following space requirements for jobs 1-4 and memory blocks. Assuming a best-fit
    scheme is used, the job,_____, is placed in the last block.
    Jobs:
    J1 10K
    J2 20K
    J3 30K
    J4 10K
    Blocks:
    B1 30K
    B2 15K
    B350K
    B4 20K
    a. J1
                                                  c. J3
                                                  d. J4
    b. J2
                         PTS: 1
    ANS: B
                                              REF: 40
13. The following algorithm can be described as _____.
     1 Set counter to 1
    2 Do while counter <= number of blocks in memory
      If job_size > memory_size(counter)
        Then counter = counter + 1
      Else
        load job into memory_size(counter)
        adjust free/busy memory lists
        go to step 4
    End do
    3 Put job in waiting queue
    4 Go fetch next job
    a. first-fit memory allocation
                                                  c. least-fit memory allocation
    b. best-fit memory allocation
                                                  d. fixed partition memory allocation
    ANS: A
                         PTS: 1
                                              REF: 41
14. The following algorithm can be described as _____.
     1 Initialize memory_block(0) = 99999
    2 Compute initial_memory_waste = memory_block(0) - job_size
    3 Initialize subscript = 0
    4 Set counter to 1
    5 Do while counter <= number of blocks in memory
         If job_size > memory_size(counter)
            Then counter = counter + 1
         Else
            memory_waste = memory_size(counter) - job_size
         If initial_memory_waste > memory_waste
            Then subscript = counter
            initial memory waste = memory waste
            counter = counter + 1
```

	Else Load job into r		•	least-fit memory allocation fixed partition memory allocation			
	ANS: B	PTS: 1	REF:	42			
15.		eginning address	of the block g Block Size 105 5 600 20 205 4050 230 1000		3		
16.	b. Memory Managc. Memory Managd. Memory Manag"free."	ger releases the bl ger immediately g ger adds block to ger resets the statu	ock and comb gives memory free list and re us of the memo	bines it with another free block. to another program. removes it from busy list. nory block where the job was stored to			
	ANS: D	PTS: 1	REF:	44			
17.	 In a dynamic partition scheme,, is how the Memory Manager deallocates a block that is betwee two other free blocks? a. The sizes of the three free partitions must be combined. b. All three are moved individually from the busy list to the free list. c. The block is combined with the larger of the two adjacent blocks. d. The status of the block is set to free. 						
	ANS: A	PTS: 1	REF:	46			
18.	When memory is de a. blank line b. null entry ANS: B	eallocated, an entr	c.	ioved from the free list by creating a(n) joined entry empty entry			
4.0							
19.	A(n) in the bu returned to the free a. blank line		-	block between two other busy memory blocks is joined entry	}		

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	b. null entry			d.	empty entry			
	ANS: B	PTS:	1	REF:	47			
20.				d section c.	bage collection or defragmentation, is performed by as of the memory space. compaction reallocation			
	ANS: C	PTS:	1	REF:	48			
21.	operation code.	ı can tel	1 the of		up of digits by its location in the line and the			
	a. functionb. value				order assignment			
	ANS: A	PTS:	1	REF:				
22.	The contains the able to access the cor a. busy list b. compaction moni	rect me		each address referenced in the program so it will be relocation. relocation register bounds register				
	ANS: C	PTS:	1	REF:	52			
23.	is the actual m a. 1,800 b. 18,000	nemory	address for a	c.	starts at 18K. 18,432 180,000			
	ANS: C	PTS:	1	REF:	53			
24.	 4. By compacting and relocating, the Memory Manager optimizes the use of memory and the throughput, but an unfortunate side effect is more a. null entries							
	ANS: D	PTS:	1	REF:	53			
25.	One approach to perform a. byte b. percentage	orming	compaction i	c.	when a certain of memory becomes busy. bit area			
	ANS: B	PTS:	1	REF:	53			